

Q2.

(i) Halting problem of TM:-

In Computability theory, the halting problem is the problem of determining from a description of the ordinary computer program and an input, whether the program will finish running or continue to run forever.

For a program β that might determine if program halt, a "pathological" program g called with an input can pass its own source and its input to β and then specially do the opposite of what β predicts g will do. No β can exist that handle this case. A proof of mathematical definition of computer and program, which became known as Turing machine.

(ii) Recursive language.

There are two equivalent major definitions for the concept of a recursive language:-

- a) A recursive formal language is a recursive subset in the set of all possible words over the alphabet of the language.
- b) A recursive language is a formal language for which there exists a Turing machine that, when presented with any finite input string, halts & accepts if the string is in the language & halts & rejects otherwise.

(iii) Variety of Turing machine

- (a) Multiple track Turing machine
- (b) Two-way infinite tape Turing machine
- (c) Multi-tape Turing machine
- (d) Multi-tape Multi-head Turing machine
- (e) Multi-directional tape Turing machine
- (f) Multi-head Turing machine
- (g) Non-deterministic Turing machine.